

# Analog to digital conversion

Sampling and reconstruction

# In the last lecture

- ▶ Non-Uniform quantizers
- ▶ Vector Quantizers
- ▶ Transformation of distributions

# In the last lecture

- ▶ Why do we consider non uniform quantizers
- ▶ Why do we look at quantizers which take multiple bits in inputs
- ▶ Why do we want to map a distributions to the uniform distribution

# In this lecture

- ▶ Mapping of quantized labels to a bit-stream
- ▶ Vector Quantizers
- ▶ Waveform coding

# Encoding

How do we map the samples to the bits

# Encoding

- ▶ As the number of quantization level increases, mapping samples to bits becomes more involved
  - ▶ Does it matter?
    - ▶ As far as we are transmitting over a noiseless channel, a one-to-one mapping is as good as another one
    - ▶ If we introduce noise, things change radically
      - ▶ We want the erroneous reconstruction to still be close to real values

# Encoding

- ▶ A simple mapping is the one in which we simply consider the binary expansion of the quantization levels
  - ▶ Level 1 001
  - ▶ Level 2 010
  - ▶ Level 3 011
  - ▶ ...

# Encoding

- ▶ What can we do if we want to minimize the impact on the reconstruction error
  - ▶ In this case we want to minimize the difference between the bit representation of one label and the next one
  - ▶ This is not the case in the natural encoding
    - ▶ 7 -> 0 1 1 1
    - ▶ 8 -> 1 0 0 0



# Encoding

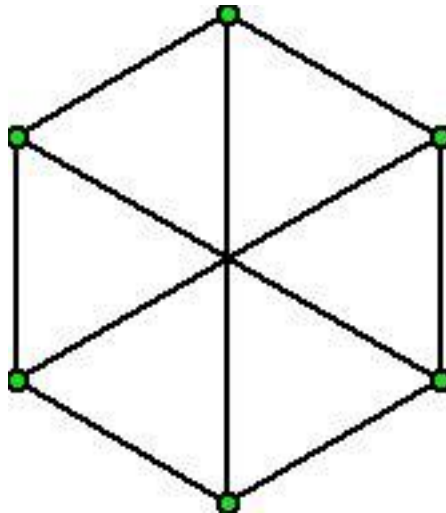
- ▶ Think about this problem:
  - ▶ This is a labelling, so we can settle for a table which maps the quantization label to a bit string

# Encoding

- ▶ Think about this problem:
  - ▶ The minimum distance between two labels is 1 bit
    - ▶ Can we attain this?
  - ▶ This at this problem in a geometrical manner

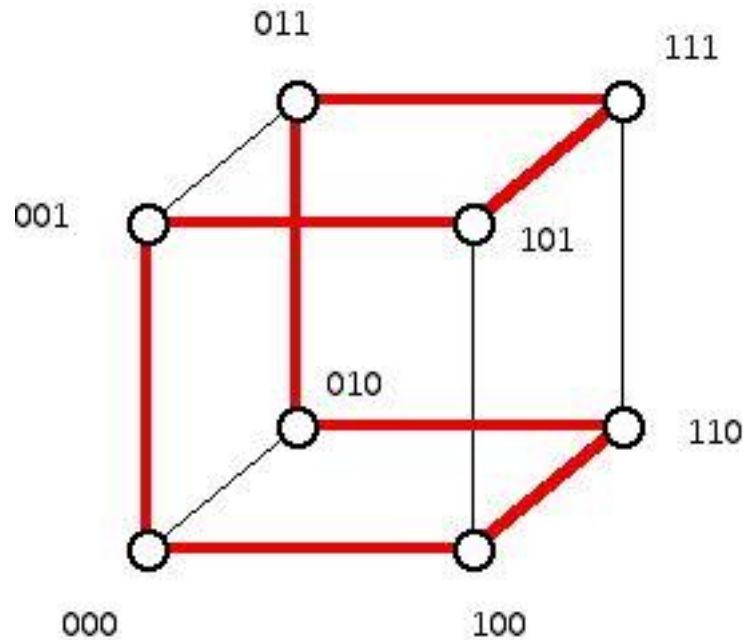
# Encoding

- ▶ A code of length 3 in a cube on 3 dimensions



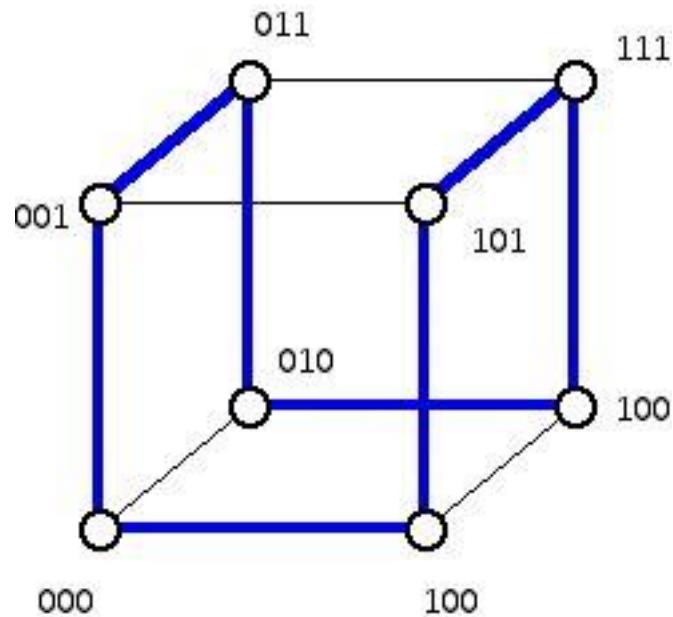
# Encoding

- ▶ The code requirement on the label distance becomes a circuit on the vertexes on the cube



# Encoding

- ▶ The code requirement on the label distance becomes a circuit on the vertexes on the cube



# Encoding

**TABLE 7.3** NBC AND GRAY CODES FOR A 16-LEVEL QUANTIZATION

Quantization level	Level order	NBC code	Gray code
$\hat{x}_1$	0	0000	0000
$\hat{x}_2$	1	0001	0010
$\hat{x}_3$	2	0010	0011
$\hat{x}_4$	3	0011	0001
$\hat{x}_5$	4	0100	0101
$\hat{x}_6$	5	0101	0100
$\hat{x}_7$	6	0110	0110
$\hat{x}_8$	7	0111	0111
$\hat{x}_9$	8	1000	1111
$\hat{x}_{10}$	9	1001	1110
$\hat{x}_{11}$	10	1010	1100
$\hat{x}_{12}$	11	1011	1101
$\hat{x}_{13}$	12	1100	1001
$\hat{x}_{14}$	13	1101	1000
$\hat{x}_{15}$	14	1110	1010
$\hat{x}_{16}$	15	1111	1011

# Waveform Coding

How is voice coded in the phone line

# Waveform Coding

- ▶ Having a statistical description of the input is not always possible
- ▶ Or sometimes the description is simply too complex or not too useful



# Waveform Coding

- ▶ How do we design a system for these scenarios
  - ▶ We can simply focus on a wide class of signals that are characterized by some reasonable features

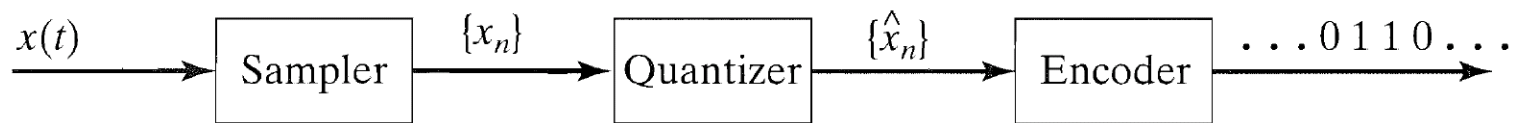
# Waveform Coding

Pulse code modulation is the simplest and oldest waveform coding scheme. A pulse code modulator consists of three basic sections: a sampler, a quantizer and an encoder. A functional block diagram of a PCM system is shown in Figure 7.7. In PCM, we make the following assumptions:

1. The waveform (signal) is bandlimited with a maximum frequency of  $W$ . Therefore, it can be fully reconstructed from samples taken at a rate of  $f_s = 2W$  or higher.
2. The signal is of finite amplitude. In other words, there exists a maximum amplitude  $x_{\max}$  such that for all  $t$ , we have  $|x(t)| \leq x_{\max}$ .
3. The quantization is done with a large number of quantization levels  $N$ , which is a power of 2 ( $N = 2^\nu$ ).

# Waveform Coding

► System model



**Figure 7.7** Block diagram of a PCM system.

# Waveform Coding

- ▶ We can start focusing on uniform sources

**Uniform PCM.** In uniform PCM, we assume that the quantizer is a uniform quantizer. Since the range of the input samples is  $[-x_{\max}, +x_{\max}]$  and the number of quantization levels is  $N$ , the length of each quantization region is given by

$$\Delta = \frac{2x_{\max}}{N} = \frac{x_{\max}}{2^{\nu-1}}. \quad (7.4.1)$$

# Waveform Coding

- ▶ What is the distortion in this case?

The distortion introduced by quantization (quantization noise) is therefore

$$E[\tilde{X}^2] = \int_{-\frac{\Delta}{2}}^{+\frac{\Delta}{2}} \frac{1}{\Delta} \tilde{x}^2 d\tilde{x} = \frac{\Delta^2}{12} = \frac{x_{\max}^2}{3N^2} = \frac{x_{\max}^2}{3 \times 4^\nu}, \quad (7.4.2)$$

where  $\nu$  is the number of bits/source sample and we have employed Equation (7.4.1). The signal-to-quantization noise ratio then becomes

$$\text{SQNR} = \frac{P_X}{\tilde{X}^2} = \frac{3 \times N^2 P_X}{x_{\max}^2} = \frac{3 \times 4^\nu P_X}{x_{\max}^2}, \quad (7.4.3)$$

# Waveform Coding

- ▶ What is the relationship between the error and the quantization levels?

Expressing SQNR in decibels, we obtain

$$\text{SQNR}|_{\text{dB}} \approx 10 \log_{10} \frac{P_X}{x_{\max}^2} + 6\nu + 4.8. \quad (7.4.4)$$

We can see that each extra bit (increase in  $\nu$  by one) increases the SQNR by 6 dB. This is a very useful strategy for estimating how many extra bits are required to achieve a desired SQNR.

# Waveform Coding

## Example 7.4.1

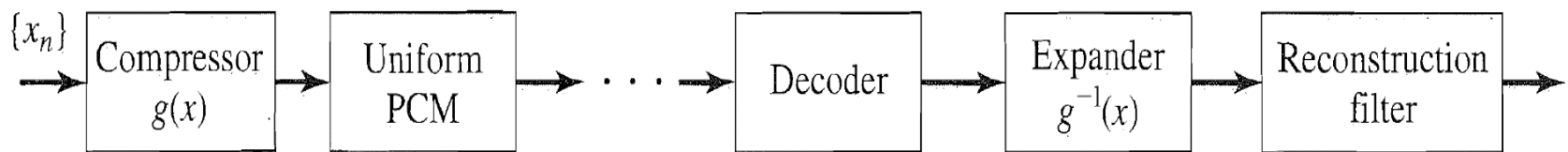
What is the resulting SQNR for a signal uniformly distributed on  $[-1, 1]$ , when uniform PCM with 256 levels is employed?

**Solution** We have  $P_X = \int_{-1}^1 \frac{1}{2}x^2 dx = \frac{1}{3}$ . Therefore, using  $x_{\max} = 1$  and  $\nu = \log 256 = 8$ , we have

$$\text{SQNR} = 3 \times 4^\nu \times P_X = 4^\nu = 65536 \approx 48.16 \text{ dB.} \quad \blacksquare$$

# Waveform Coding

The usual method for performing nonuniform quantization<sup>3</sup> is to first pass the samples through a nonlinear element that compresses the large amplitudes (reduces the dynamic range of the signal) and then performs a uniform quantization on the output. At the receiving end, the inverse (expansion) of this nonlinear operation is applied to obtain the sampled value. This technique is called *companding* (*compressing–expanding*). A block diagram of this system is shown in Figure 7.8.



**Figure 7.8** Block diagram of a nonuniform PCM system.



# Waveform Coding

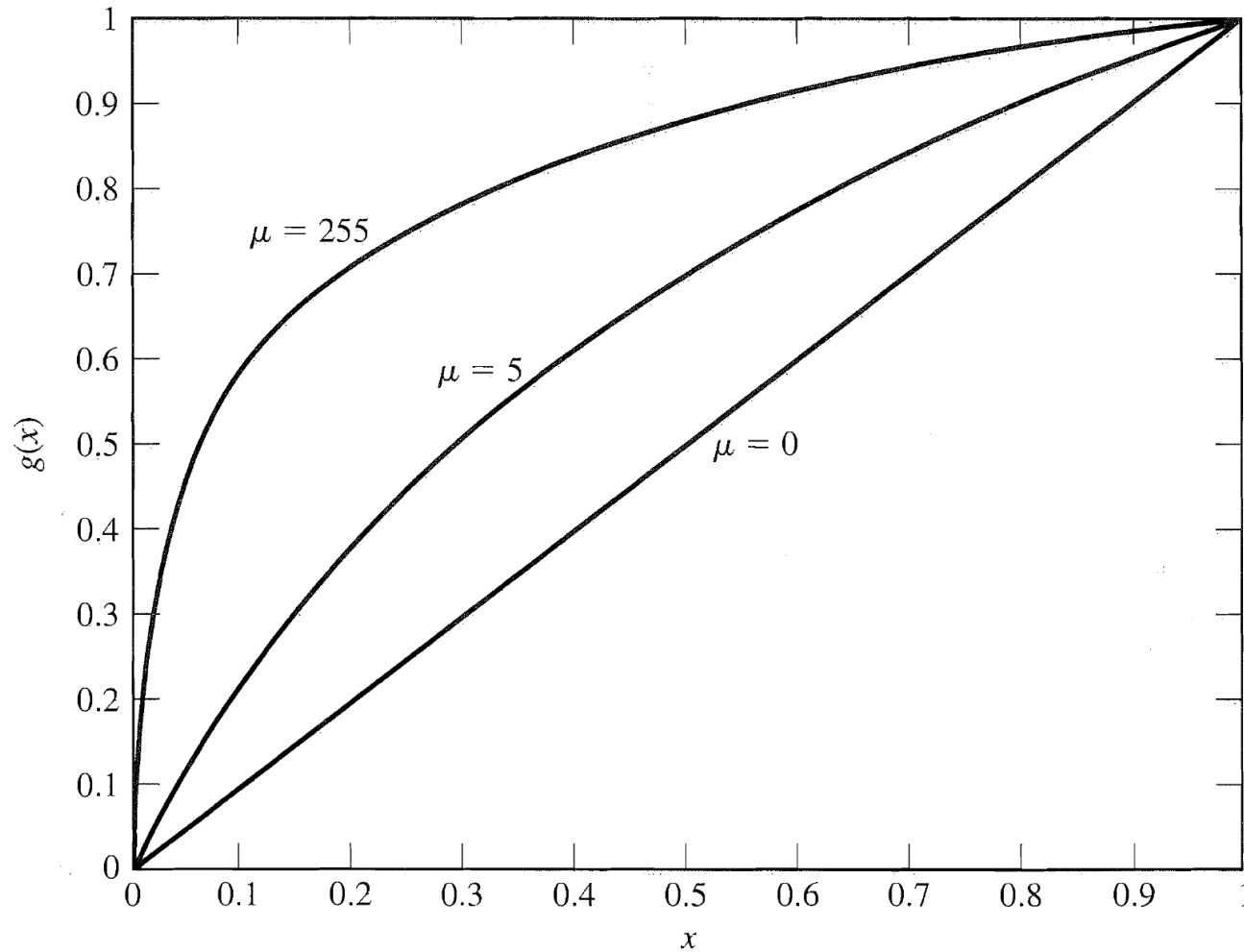
- ▶ In telephone systems we often use a pre/post processing of the information

There are two types of companders that are widely used for speech coding. The  $\mu$ -law compander, used in the United States and Canada, employs the logarithmic function at the transmitting side, where  $|x| \leq 1$ :

$$g(x) = \frac{\log(1 + \mu|x|)}{\log(1 + \mu)} \operatorname{sgn}(x). \quad (7.4.7)$$

The parameter  $\mu$  controls the amount of compression and expansion. The standard PCM system in the United States and Canada employs a compressor with  $\mu = 255$  followed by a uniform quantizer with 8 bits/sample. Use of a compander in this system improves the performance of the system by about 24 dB. Figure 7.9 illustrates the  $\mu$ -law compander characteristics for  $\mu = 0, 5, \text{ and } 255$ .

# Waveform Coding



**Figure 7.9** A graph of  $\mu$ -law compander characteristics.

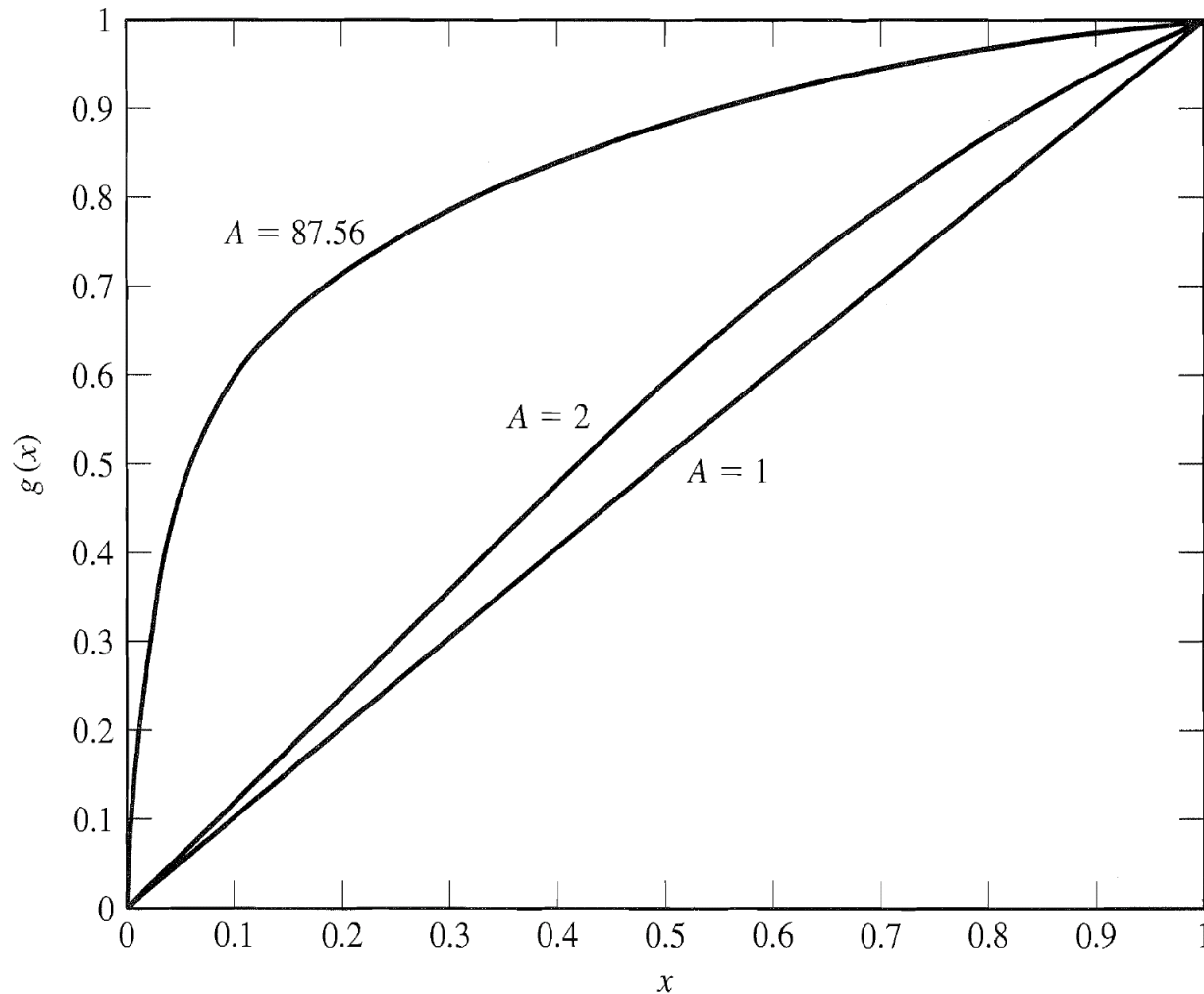
# Waveform Coding

The second widely used logarithmic compressor is the  $A$ -law compander. The characteristics of this compander are given by

$$g(x) = \frac{1 + \log A|x|}{1 + \log A} \operatorname{sgn}(x), \quad (7.4.8)$$

where  $A$  is chosen to be 87.56. The performance of this compander is comparable to the performance of the  $\mu$ -law compander. Figure 7.10 illustrates the characteristics of this compander for  $A = 1, 2,$  and 87.56.

# Waveform Coding



**Figure 7.10** A graph of A-law compander characteristics.

The background features abstract, overlapping geometric shapes in various shades of blue, ranging from light to dark, on the right side of the slide. The rest of the background is white.

**That's it for today!**

I'll see you next time